

Lecture 2: Simultaneous Move Games Theory

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In the Previous Lecture

- Game Theory: set of tools to analyse behaviour in the presence of strategic interdependence
 - For example, firms in an oligopoly market
 - Static games: players play simultaneously and only once
 - Further examples: "prisoner's dilemma", "coordination game", "matching pennies", "stag hunt", ...
 - Representation in "normal form"
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Today's Lecture

- What should we expect players to play?
 - Looking for reasonable concepts in simple predictable games and apply these concepts in other settings
 - In this chapter, concentrate in static ("simultaneous move" or "strategic") games
 - Solution concepts:
 - Use dominant strategies
 - Don't use dominated strategies
 - Play Nash equilibrium strategies
 - Assume that players are rational:
 - each chooses her best action according to her preferences
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Dominant Strategies

- "Prisoner's dilemma":

1\2	<i>DC</i>	<i>C</i>
<i>DC</i>	-2,-2	-10,-1
<i>C</i>	-1,-10	-5,-5

- What action would you choose?
 - No matter what the other does, it is better to play "Confess":
If she plays *DC*, one obtains -1 by playing *C* and -2 by playing *DC*
If she plays *C*, one obtains -5 by playing *C* and -10 by playing *DC*
 - Formally, "Confess" strictly dominates "Don't Confess"
 - "Confess" is a dominant strategy or action
 - Both confessing is the outcome! Conflict with Pareto-optimality
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Introducing Notation

- $i = 1, \dots, I$ the players of a game (e.g. in PD $i = 1, 2$)
 - s_i a strategy or action for player i (e.g. $s_1 = C$ or $s_1 = DC$)
 - $s = (s_1, \dots, s_I)$ a strategy profile: one strategy for each player (e.g. $s = (C, DC)$ or $s = (CD, DC)$)
 - $s = (s_i, s_{-i})$ again a strategy profile, where $-i$ means all the players except i
 - $u_i(s) = u_i(s_1, \dots, s_I)$ the payoff for player i if s is played (e.g. $u_1(C, DC) = -1$, $u_2(C, DC) = -10, \dots$)
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Dominant Strategies

- Definition: Player i 's strategy s_i'' strictly dominates strategy s_i' if
$$u_i(s_i'', s_{-i}) > u_i(s_i', s_{-i})$$
 for every possible list s_{-i} of the other player's actions
 - Accordingly, we say that strategy s_i' is strictly dominated
 - Example: strategy "Don't Confess" is strictly dominated by strategy "Confess"
 - Definition: A strictly dominant strategy for player i is a strategy that strictly dominates all the other strategies
 - Example: strategy "Don't Confess" is strictly dominant for both players
 - Rational players play dominant strategies
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Strictly Dominated Strategies

- Problem: strictly dominant strategies rarely exist. Examples:

(a)

$1 \backslash 2$	L	R
U	1,-1	-1,-1
M	-1,1	1,-1
D	-2,5	-3,2

(b)

$1 \backslash 2$	L	R
U	5,1	4,0
M	6,0	3,1
D	6,4	4,4

- However, a strictly dominated strategy may still exist.
 - Example: strategy D in game (a)
 - Rational players do not play dominated strategies
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Weakly Dominated Strategies

- Definition: Player i 's strategy s_i'' weakly dominates strategy s_i' if

$u_i(s_i'', s_{-i}) \geq u_i(s_i', s_{-i})$ for every possible list s_{-i} of other player's actions
with strict inequality for some s_{-i}

- Example: U and M in (b) are weakly dominated
 - Should we rule out weakly dominated strategies as well?
 - No! Playing M can be as good as playing D if 1 *believes* that 2 will play L
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Nash Equilibrium

- Problem: strictly dominated strategies may not exist. Example: BoS game

1\2	Bach	Stravinsky
Bach	2,1	0,0
Stravinsky	0,0	1,2

- Definition: A strategy profile (s_1, s_2, \dots, s_I) constitutes a Nash equilibrium if for every player i

$$u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i}) \text{ for every action } s'_i$$

- Here, we are assuming that...
 - (1) Players are rational (given the belief about others' actions)
 - (2) Their beliefs about the actions of the others are correct
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Example 1: Prisoner's Dilemma

1\2	Don't Confess	Confess
Don't Confess	-2,-2	-10,-1
Confess	-1,-10	-5,-5

- (C, C) is a NE:
given that 2 plays C, then playing C is better than DC for 1 (-5 instead of -10)
given that 1 plays C, then playing C is better than DC for 2 (-5 instead of -10)
 - (C, DC) is a not a NE:
given that 1 plays C, then playing C is better than DC for 2 (-5 instead of -10)
 - (DC, C) is a not a NE:
given that 2 plays C, then playing C is better than DC for 1 (-5 instead of -10)
 - (DC, DC) is a not a NE:
given that 2 plays DC, then playing C is better than DC for 1 (-1 instead of -2)
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Remarks

1. NE may not be unique!! Example 2: coordination Game (or BoS)

1\2	Bach	Stravinsky
Bach	2,1	0,0
Stravinsky	0,0	1,2

- NE: (Bach, Bach) and (Stravinsky, Stravinsky)

2. NE may not exist! Example 3: Matching Pennies

1\2	Head	Tail
Head	1,-1	-1,1
Tail	-1,1	1,-1

- NE: none!
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Exercise

- Exercise: what are the NE in the Stag Hunt game (example 4)?

1\2	Stag	Hare
Stag	2,2	0,1
Hare	1,0	1,1

- Find them intuitively first and show it formally afterwards (using the previous definition)
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