

# Lecture 2: Predicting outcomes

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## In the Previous Chapter

- Game Theory: set of tools to analyse behaviour in the presence of strategic interdependence
  - For example, firms in an oligopoly market
  - Static games: players play simultaneously and only once
  - Further examples: "prisoner's dilemma", "battle of the sexes", "matching pennies", "stag hunt", ...
  - Representation in "normal form"
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## This chapter's Plan

- What should we expect players to play?
  - Looking for reasonable concepts in simple predictable games and apply these concepts in other settings
  - In this chapter, concentrate in static ("simultaneous move" or "strategic") games
  - Solution concepts:
    - Use dominant strategies
    - Don't use dominated strategies
    - Play Nash equilibrium strategies
  - Assume that players are rational:
    - each chooses her best action according to her preferences
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## Dominant Strategies

- "Prisoner's dilemma":

|           |           |          |
|-----------|-----------|----------|
| 1\2       | <i>DC</i> | <i>C</i> |
| <i>DC</i> | -2,-2     | -10,-1   |
| <i>C</i>  | -1,-10    | -5,-5    |

- What action would you choose?
  - No matter what the other does, it is better to play "Confess":  
If she plays *DC*, one obtains  $-1$  by playing *C* and  $-2$  by playing *DC*  
If she plays *C*, one obtains  $-5$  by playing *C* and  $-10$  by playing *DC*
  - Formally, "Confess" strictly dominates "Don't Confess"
  - "Confess" is a dominant strategy or action
  - Both confessing is the outcome! Conflict with Pareto-optimality
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## Introducing Notation

- $i = 1, \dots, I$  the players of a game (e.g. in PD  $i = 1, 2$ )
  - $s_i$  a strategy or action for player  $i$  (e.g.  $s_1 = C$  or  $s_1 = DC$ )
  - $s = (s_1, \dots, s_I)$  a strategy profile: one strategy for each player (e.g.  $s = (C, DC)$  or  $s = (CD, DC)$ )
  - $s = (s_i, s_{-i})$  again a strategy profile, where  $-i$  means all the players except  $i$
  - $u_i(s) = u_i(s_1, \dots, s_I)$  the payoff for player  $i$  if  $s$  is played (e.g.  $u_1(C, DC) = -1$ ,  $u_2(C, DC) = -10, \dots$ )
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## Dominant Strategies

- Definition: Player  $i$ 's strategy  $s_i''$  strictly dominates strategy  $s_i'$  if  $u_i(s_i'', s_{-i}) > u_i(s_i', s_{-i})$  for every possible list  $s_{-i}$  of the other player's actions
  - Accordingly, we say that strategy  $s_i'$  is strictly dominated
  - Example: strategy "Don't Confess" is strictly dominated by strategy "Confess"
  - Definition: A strictly dominant strategy for player  $i$  is a strategy that strictly dominates all her other strategies
  - Example: strategy "Don't Confess" is strictly dominant for both players
  - Rational players play dominant strategies
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## Strictly Dominated Strategies

- Problem: strictly dominant strategies rarely exist. Examples:

(a)

| $1 \backslash 2$ | $L$  | $R$   |
|------------------|------|-------|
| $U$              | 1,-1 | -1,-1 |
| $M$              | -1,1 | 1,-1  |
| $D$              | -2,5 | -3,2  |

(b)

| $1 \backslash 2$ | $L$ | $R$ |
|------------------|-----|-----|
| $U$              | 5,1 | 4,0 |
| $M$              | 6,0 | 3,1 |
| $D$              | 6,4 | 4,4 |

- However, a strictly dominated strategy may still exist.
  - Example: strategy D in game (a)
  - Rational players do not play dominated strategies
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## Weakly Dominated Strategies

- Definition: Player  $i$ 's strategy  $s_i''$  weakly dominates strategy  $s_i'$  if

$u_i(s_i'', s_{-i}) \geq u_i(s_i', s_{-i})$  for every possible list  $s_{-i}$  of other player's actions  
with strict inequality for some  $s_{-i}$

- Example: U and M in (b) are weakly dominated
  - Should we rule out weakly dominated strategies as well?
  - No! Playing M can be as good as playing D if 1 *believes* that 2 will play L
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## Nash Equilibrium

- Problem: strictly dominated strategies may not exist. Example: BoS game

|            |      |            |
|------------|------|------------|
| 1\2        | Bach | Stravinsky |
| Bach       | 2,1  | 0,0        |
| Stravinsky | 0,0  | 1,2        |

- Definition: A strategy profile  $(s_1, s_2, \dots, s_I)$  constitutes a Nash equilibrium if for every player  $i$

$$u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i}) \text{ for every action } s'_i$$

- Here, we are assuming that...
    - (1) Players are rational (given the belief about others' actions)
    - (2) Their beliefs about the actions of the others are correct
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## Example 1: Prisoner's Dilemma

|               |               |         |
|---------------|---------------|---------|
| 1\2           | Don't Confess | Confess |
| Don't Confess | -2,-2         | -10,-1  |
| Confess       | -1,-10        | -5,-5   |

- (C, C) is a NE:  
given that 2 plays C, then playing C is better than DC for 1 (-5 instead of -10)  
given that 1 plays C, then playing C is better than DC for 2 (-5 instead of -10)
  - (C, DC) is a not a NE:  
given that 1 plays C, then playing C is better than DC for 2 (-5 instead of -10)
  - (DC, C) is a not a NE:  
given that 2 plays C, then playing C is better than DC for 1 (-5 instead of -10)
  - (DC, DC) is a not a NE:  
given that 2 plays DC, then playing C is better than DC for 1 (-1 instead of -2)
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## Remarks

1. NE may not be unique!! Example 2: coordination Game (or BoS)

|            |      |            |
|------------|------|------------|
| 1\2        | Bach | Stravinsky |
| Bach       | 2,1  | 0,0        |
| Stravinsky | 0,0  | 1,2        |

- NE: (Bach, Bach) and (Stravinsky, Stravinsky)

2. NE may not exist! Example 3: Matching Pennies

|      |      |      |
|------|------|------|
| 1\2  | Head | Tail |
| Head | 1,-1 | -1,1 |
| Tail | -1,1 | 1,-1 |

- NE: none!
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## Exercise

- Exercise: what are the NE in the Stag Hunt game (example 4)?

|      |      |      |
|------|------|------|
| 1\2  | Stag | Hare |
| Stag | 2,2  | 0,1  |
| Hare | 1,0  | 1,1  |

- Find them intuitively first and show it formally afterwards (using the previous definition)
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